# Responsive and Flexible CNL Authoring with Zipper-based Transformations



Responsive and Flexible CNL Authoring

CNL'18 1/24

#### **Overview**

- 1) The FL-NL Language Gap
- Principles of the N<A>F Design Pattern
- Illustration on a Core RDF Query Language
- 4 Application to 3 Semantic Web Tasks
- 5 Conclusion

# The Gap between Formal Languages and Natural Languages



• Humans speak English, French, Chinese, ...

Natural Languages (NL)

Machines speak RDF, OWL, SPARQL, ...

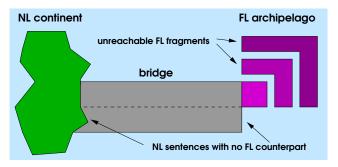
Formal Languages (FL)

Only a few humans speak both...
 ... so we need bridges over the gap

#### The Problem of Adequacy

Adequacy = expressivity + safeness

- an essential property of language bridges
- expressivity (~recall): proportion of FL sentences reachable through the bridge
- safeness (~precision): proportion of paths on the bridge leading to correct FL sentences



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#### **Different Kinds of Bridges**

Different approaches have been explored to cross the gap for search:

- Question Answering (QA): "unsafe full-way bridge"
  - users express questions in spontaneous NL
  - systems often fail to understand the question or cannot answer it
  - Iow coverage of target FL
- Controlled Natural Languages (CNL): "safe half-way bridge"
  - wide coverage of target FL
  - users must use restricted grammar and lexicon
  - systems can help write well-formed questions (autocompletion)
- Query Builders (QB): "safe and assisted climbing"
  - users still have to build formal queries
  - systems help build well-formed queries

They offer different trade-offs between expressivity (FL coverage), safeness (reliability), and readability (closeness to NL).

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# Limits of Autocompletion for CNL Authoring

#### Autocompletion

Suggest the next possible words according to the grammar and lexicon.

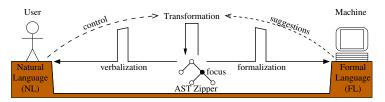
#### Limits

- Responsiveness
  - partial sentence at most steps
  - hence no translation/interpretation in FL
  - hence lack of feedback: e.g. query results
- Flexibility
  - in general, only adding words at the end
  - sometimes, one word at a time
  - restricted edition compared to text editors (cursor)

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# Bridging the Gap with Zippers

The N<A>F design pattern for responsive and flexible CNL authoring.



A kind of "suspended bridge":

pillar Abstract Syntax Trees (AST) + Huet's zippers for focus suspender transformations of AST zippers decks translations defined as Montague grammars cables system suggestions and user control

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#### **Pros and Cons**

#### PROS

- bridges the NL-FL gap because two-way synchronous translations
- scales in expressivity because ambiguities are solved piecewise during building
- ensures strong safeness because fine-grained guidance during building
- Is responsive because a complete sentence is defined at all steps
- offers a lot of flexibility because building applies to a tree, not a sequence of words, and focus as cursor
- applies to various tasks because no assumption is made on the FL CONS
  - does not apply to spontaneous NL or existing texts
  - As slower interaction because of the building process

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#### Illustration on a Core RDF Query Language (CRQL)

To show a concrete application of the N < A > F design pattern

- task: semantic search on RDF data
- formal language: CRQL, a fragment of SPARQL tree patterns, unions, negations
- safeness criteria: avoid empty results

Bridge components:

- ASTs
- AST zippers for focus representation
- In the second second
- Itranslation to SPARQL
- translation to English
- Computation of suggestions

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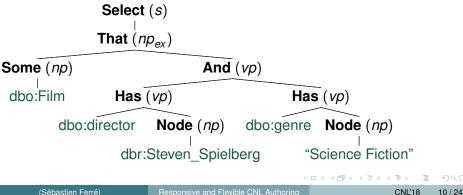
- ASTs
- AST zippers for focus representation
- ST zipper transformations for AST building
- translation to SPARQL
- Itranslation to English
- 6 computation of suggestions

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#### 1. CRQL ASTs

ASTs are close to NL syntax but much more abstract

- sentences (s) denote queries
- noun phrases (np) denote sets of entities
- verb phrases (vp) denote conditions on entities
- words are RDF classes, properties, and nodes



#### 1. ASTs Specification

ASTs are trees that can be specified with algebraic datatypes\*:

$$s :=$$
**Select**( $np$ )

det := Some | Every | No

\* source code online in ML style (OCaml)

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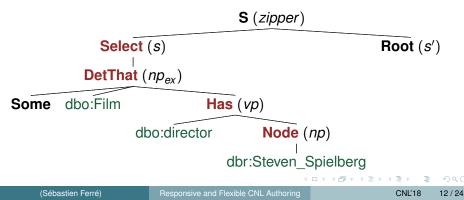
Huet's Zippers (functional pearl at J. Functional Prog., 1997)

- type-safe representation of focus in complex data structures
- efficient focus-centered edition of data structures (transformations)
- open and close data structures like a jacket!
- s', np', vp' are datatypes for the contexts of s, np, vp
- zipper = sub-tree under focus + context

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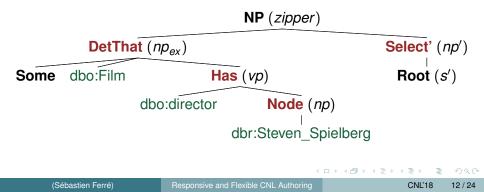
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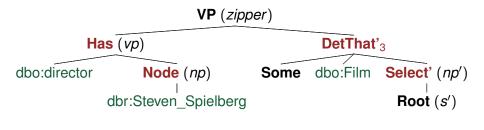
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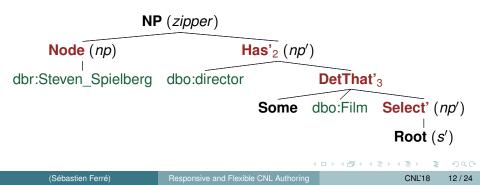
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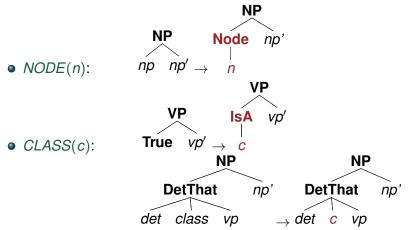
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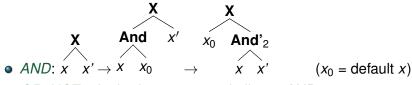
#### 3. AST Zipper Transformations (1/3)

A transformation goes from zipper to zipper, used as a building step



• DET(d), PROP(p): insertions, similar to NODE and CLASS

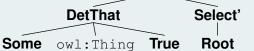
3. AST Zipper Transformations (2/3)



- OR, NOT: algebraic operators, similar to AND
- DOWN, UP, LEFT, RIGHT: focus moves

#### Theorem

The set of transformations is CRQL-complete from initial zipper



#### 3. AST Zipper Transformations (3/3)

The above example AST is reached by the following sequence of transformations:

- CLASS(dbo:film): NP(DetThat(Some,dbo:Film,True), Select'(Root))
- OWN: VP(True, DetThat'<sub>3</sub>(Some,dbo:Film,Select'(Root)))
- PROP(dbo:director): VP(Has(dbo:director,DetThat(Some,owl:Thing,True)), DetThat'<sub>3</sub>(Some,dbo:Film,Select'(Root)))
- DOWN: NP(DetThat(Some,owl:Thing,True), Has'<sub>2</sub>(dbo:director, DetThat'<sub>3</sub>(Some,dbo:Film,Select'(Root)))
- NODE(dbr:Steven\_Spielberg): NP(Node(dbr:Steven\_Spielberg), Has'<sub>2</sub>(dbo:director, DetThat'<sub>3</sub>(Some,dbo:Film,Select'(Root))))
- UP, UP, ...:

S(Select(DetThat(Some,dbo:Film,Has(dbo:director,Node(dbr:Steven\_Spielberg)))) Root)

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### 4. Translation to SPARQL (Formalization)

R. Montague's Grammar ("English as a formal language", 1970)

- designed for translation from NL to logic
- compositional semantics based on lambda calculus
- Montague grammar = grammar rules + lambda-terms
  - here, AST datatypes play the role of grammars

#### Exerpt

 $\begin{array}{ll} vp := \mathsf{IsA}(\mathit{class}) & \lambda x.(x + 'a' + \mathit{class}) \\ vp := \mathsf{Not}(vp_1) & \lambda x.('\texttt{FILTER NOT EXISTS } \{' + (vp_1 x) + '\}') \end{array}$ 

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# 4. Full Montague Grammar for Formalization in SPARQL

		<b>a i</b> i ( )	
S	:=	Select(np)	'SELECT ?x <sub>1</sub> WHERE { ' + ( $np \lambda x.(")$ ) + ' }
np	:=	Node(node)	$\lambda d.((d node))$
•		DetThat(det, cl, vp)	$\lambda d.(\det cl \lambda x.((d x) + ' \cdot ' + (vp x)))$
		$And(np_1, np_2)$	$\lambda d.((np_1 d) + ' . ' + (np_2 d))$
	Ì	$Or(np_1, np_2)$	$\lambda d.(' \{ ' + (np_1 d) + ' \}$ UNION $\{ ' + (np_2 d) + ' \}$
	Í	Not(np)	$\lambda d.(' \text{NOT} \{' + (np d) + ' \}')$
det	:=	Some	$\lambda d_1 \cdot \lambda d_2 \cdot ((d_1 ' ? x_i' + ' \cdot ' + (d_2 ' ? x_i')))$
		No	$\lambda d_1 \cdot \lambda d_2 \cdot (' \text{NOT} \{ ' + (d_1 ' ? x_i' + ' \cdot ' + (d_2 ' ? x_i') \}$
		Every	$\lambda d_1 \cdot \lambda d_2 \cdot (' \text{ NOT } \{' + (d_1 ' ? x_i' + ' \cdot \text{ NOT} \{' + (d_2 ) \} \}$
vp	:=	IsA(class)	$\lambda x.(x + 'a' + class)$
		Has(prop, np)	$\lambda x.((np \ \lambda y.(x + prop + y)))$
		IsOf(prop, np)	$\lambda x.((np \ \lambda y.(y + prop + x)))$
		True	$\lambda x.(")$
	Í	$And(vp_1, vp_2)$	$\lambda x.((vp_1 x) + ' \cdot ' + (vp_2 x))$
	İ	$Or(vp_1, vp_2)$	$\lambda x.(' \{ ' + (vp_1 x) + ' \}$ UNION $\{ ' + (vp_2 x) + ' \}$
	İ	Not(vp)	$\lambda x.(' \text{ NOT } \{' + (vp x) + ' \}')$

# 5. Translation to English (Verbalization)

Montague grammars can also be used here

- English as target language
- compositional generation of NL phrases
  - s → sentences, np → noun phrases
  - $vp \rightarrow$  relative clauses parametrized by negation ( $\lambda n$ .)
  - ► class, prop ~→ noun
- linearization in Grammatical Framework

#### Excerpt

*is* 0 = 'is'

is 1 = 'is not'

#### 4 & 5. Translation Example

The example AST above has the following translations.

# SPARQL SELECT ?x1 WHERE { ?x1 a dbo:Film . ?x1 dbo:director dbr:Steven\_Spielberg . }

#### English

Give me a film whose director is Steven Spielberg

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# 6. Computation of System Suggestions

No general technique for this component:

- depends on the task
- depends on the FL semantics
- depends on the safeness criteria

For semantic search with CRQL, suggested *insertion* transformations are computed from SPARQL results and from the focus entity  $\mathbf{x}$ 

- nodes: values of x
- classes: values of ?c s.t. { x a ?c }
- properties: values of ?p s.t. { x ?p [] } or { [] ?p x }

#### Theorem

The suggestions prevent empty results (safeness), and yet are complete w.r.t. non-empty CRQL queries (expressivity) ⇒ perfect adequacy to CRQL.

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#### Application to 3 Semantic Web Tasks

To show the effectiveness and genericity of the N<A>F design pattern

	SPARKLIS	SEWELIS/UTILIS	PEW
task	querying RDF	authoring RDF	completing OWL
	endpoints	descriptions	ontologies
FL	SPARQL	RDF	OWL
expres-	CRQL	conjunctive sub-	CRQL
sivity	+ cyclic patterns	set of CRQL	- non-atomic neg-
	+ OPTIONAL		ations
	+ ordering		
	+ expressions		
safe-	no empty results	similarity to previ-	no inconsistency
ness		ous descriptions	
sugges-	SPARQL eval.	query relaxation	satisfiability
tions			checks

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#### **Results from User Studies**

- SPARKLIS has 200-2000 hits per month since Spring 2014
- SPARKLIS has been adopted by two French institutions
- UTILIS' fine-grained suggestions prefered to Protégé's
- UTILIS help to produce more consistent data
- PEW better in quantity and quality than Protégé
  - 56% vs 24% completion in formalization of hand anatomy
  - more axioms produced with smaller error rate
- Main difficulty: understand the focus, its impact on suggestions, and the need to move it

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#### Conclusion

The N < A > F design patter is

- a powerful strategy to build bridges over the NL-FL gap
  - users are never exposed to FL (readability)
  - and machines are never exposed to NL
  - users cannot fall in the gap (safeness)
  - large subsets of FL are reachable by users (expressivity)
- an interesting alternative to CNL Autocompletion
  - formal interpretation (e.g. results) is available at all steps (responsiveness)
  - query elements can be inserted/deleted at any focus (flexibility)
  - edition steps are more semantic
     e.g. inserting a property means crossing a relation in the RDF graph
- an interesting alternative to Question Answering
  - that avoids the hard problem of NL understanding
  - that scales in expressivity in a modular way
  - that applies to various tasks and FL

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#### The End

Questions ?

(Sébastien Ferré)

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